

Representations of objects

Quadtree, octree and BSP

Meshes, half-edge and winged edge data structures

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Motivation, needs

- CAGD, displaying data, game design, entertainment
- different needs - different representations

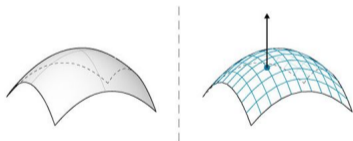


Figure: "nice" image

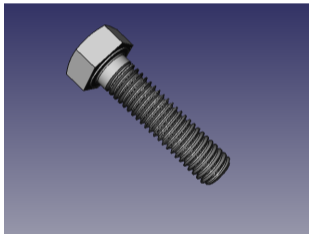


Figure: precise model

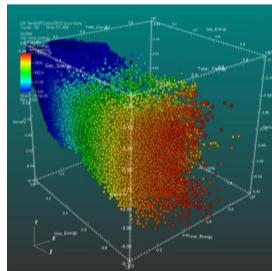
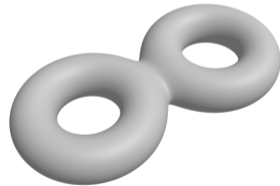
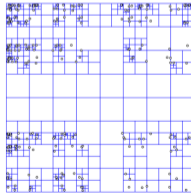
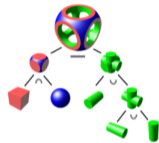
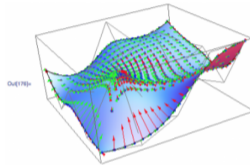
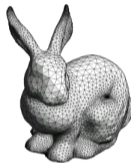


Figure: data visualization

Types of object representations

- meshes
- splines
- CSG
- space partitioning
- implicit surfaces
- etc.



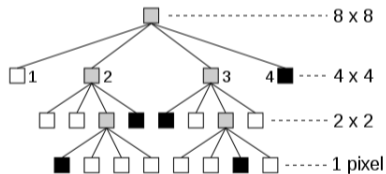
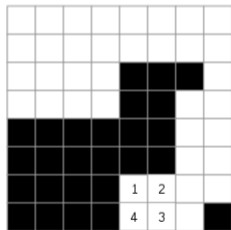
Space partitioning

- dividing (usually Euclidean) space into two or more disjoint subsets
- usage in ray tracing (wait for course CG(2))

- quadtree
- octtree
- BSP tree
- ~~k -d tree~~

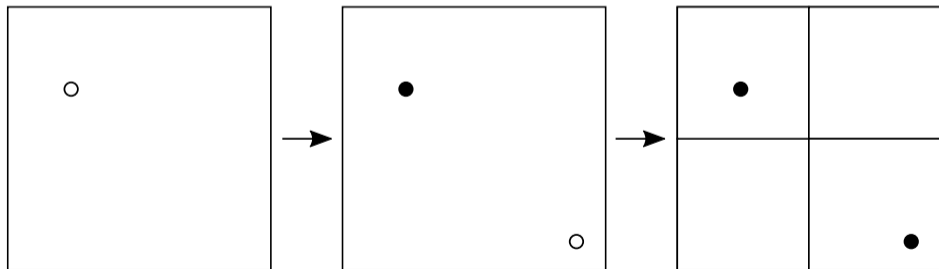
Quadtree – definition

- tree (data structure) with each branch having exactly 4 children



Quadtree – features

- decomposition of a space into adaptive cells (usually squares or rectangles)
- each cell has maximum capacity
- if the capacity has been exceeded, the cell is subdivided



Quadtree – types

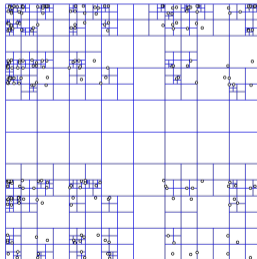
■ region quadtree

- max. capacity – one (or more) colour in a cell
- e.g. *image representation*



■ point quadtree

- max. capacity – one (or more) point in a cell
- e.g. *mesh generation*

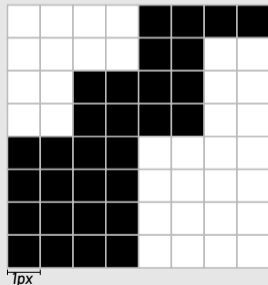


Quadtree – representations

- as a **tree** data structure
- as a **text string** (sometimes called *linear quadtree*)

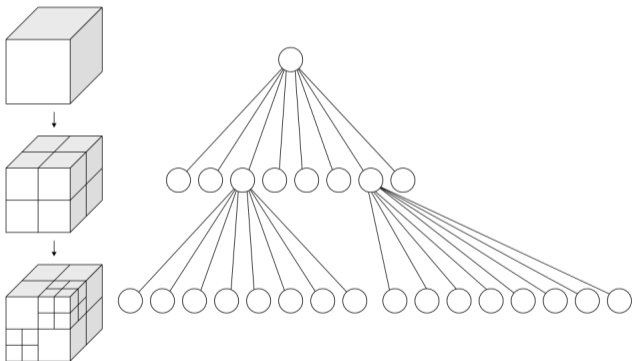
Exercise

Record the following monochromatic image using quadtree in both tree and line notation:



Octree

- each branch has exactly 8 children
- one subdivision step creates 8 octants
- 3D analogy of quadtree
- e.g. *collision detection in game design*



Binary space partitioning – BSP

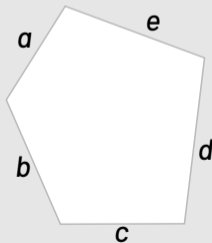
- binary tree
- one subdivision step divides a space into convex sets by **hyperplanes**
- generalization of both quadtree and octree
- store and retrieve spatial information about primitives
- static geometry
- e.g. *painter's algorithm*, *CSG - operations with shapes*

BSP – generation

- scene is divided into two until partitioning satisfies one or more criteria
- for points (???), for polygons (???)

Exercise

*Generate BSP tree for a pentagon.
What is the optimal hyperplane for
the initial subdivision?*



Vertex, edge and face

Definition

Let $\mathcal{V} = \{V_1, \dots, V_n\}$, $V_i \in \mathbb{R}^3$, $i = 1, \dots, n$. Then V_i is called a **vertex** and \mathcal{V} is called a **set of vertices**.

Definition

Denote $\mathcal{E} \subseteq \mathcal{V}^{[2]}$, where $\mathcal{V}^{[2]} = \{[P_i, P_j] \in \mathcal{V}^2 \mid P_i \neq P_j\}$ and each $e = [V_i, V_j] \in \mathcal{E}$ satisfies $[V_j, V_i] = e$. Then e is called an **edge** and \mathcal{E} is called a **set of edges**.

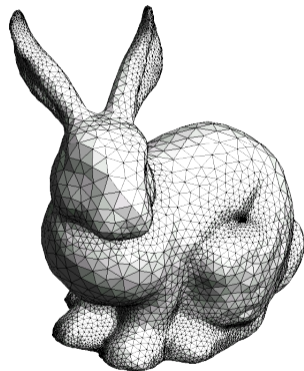
Definition

Let $\mathcal{F} \subseteq \mathcal{V}^{[n]}$, $n \geq 3$, where $\mathcal{V}^{[n]} = \{[V_{i_1}, V_{i_2}, \dots, V_{i_n}] \in \mathcal{V}^n \mid V_{i_j} \neq V_{i_k}, j \neq k\}$. The set \mathcal{F} is called a **set of faces**. Its element $F = [V_{i_1}, V_{i_2}, \dots, V_{i_n}] \in \mathcal{F}$ is called a **face** and conforms the condition, that all permutations of $V_{i_1}, V_{i_2}, \dots, V_{i_n}$ represent the same face F .

Definition

Denote $\mathcal{M} = \{\mathcal{V}, \mathcal{E}, \mathcal{F}\}$. \mathcal{M} is called a **polygon mesh**. The corresponding sets of vertices, edges and faces are denoted by $\mathcal{V}(\mathcal{M})$, $\mathcal{E}(\mathcal{M})$ and $\mathcal{F}(\mathcal{M})$, respectively.

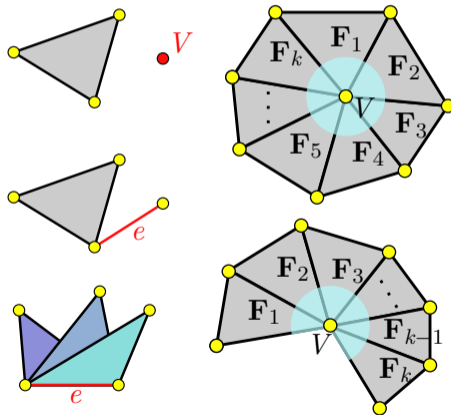
A mesh is meant to be a **2-manifold**, meaning each vertex on a mesh \mathcal{M} has a neighbourhood that is homeomorphic to the Euclidean space \mathbb{E}^2 .



Mesh

To avoid non-manifold meshes, additional conditions are required:

- 1 each vertex $V \in \mathcal{V}$ lies on some of the edges from \mathcal{E}
- 2 each edge $e \in \mathcal{E}$ lies on some of the faces from \mathcal{F}
- 3 each edge $e \in \mathcal{E}$ belongs to at most two faces from \mathcal{F}
- 4 for each face $[P_i, P_j, P_k] \in \mathcal{F}$ we have $[P_i, P_j], [P_i, P_k], [P_j, P_k] \in \mathcal{E}$,
- 5 for each vertex $V \in \mathcal{V}$ and all faces $\mathbf{F}_1, \dots, \mathbf{F}_k \in \mathcal{F}$ sharing V , one of the following statements applies:
 - $\mathbf{F}_1, \dots, \mathbf{F}_k$ create a simple cycle around V
 - $\mathbf{F}_1, \dots, \mathbf{F}_k$ create a simple sequence around V



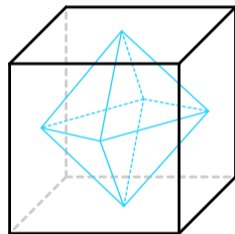
Mesh

Definition

Boundary of a mesh is the union of all edges belonging to exactly one face. The mesh is said to be **closed**, if it has no boundary.

Definition

The mesh is **connected** if the dual graph is connected.



Polygon mesh – properties

■ Euler's formula

$$\#V - \#E + \#F = 2(1 - g),$$

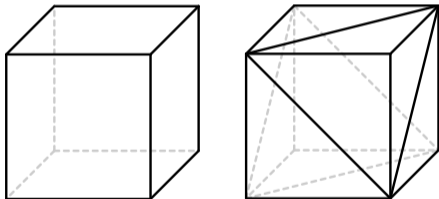
g - genus (number of handles), $\#V, \#E, \#F$ - number of vertices, edges and faces, respectively

For triangle meshes:

■ number of vertices, faces and edges

$$\#F \approx 2\#V, \quad \#E \approx 3\#V$$

■ average **vertex valence** (number of incident edges) is 6



How to represent a mesh?

- the efficiency and memory consumption depends on the data structure
- we need to take into account
 - **topological requirements** – *2-manifold meshes/complex edges or singular vertices, triangle/polygon meshes etc.*
 - **algorithmic requirements** – *just render the mesh, local neighbourhoods of vertices/edges/faces, static/changing geometry etc.*

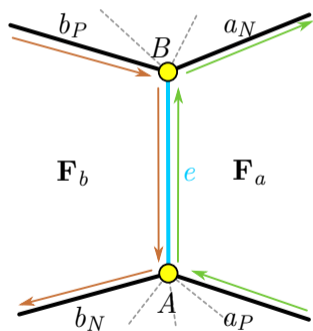
Face-based data structures

Vertices
x_1, y_1, z_1
...
x_v, y_v, z_v

Triangles
i_{11}, i_{12}, i_{13}
...
...
...
i_{F1}, i_{F2}, i_{F3}

- vertices \longleftrightarrow coordinates
- triangles \longleftrightarrow indices
- static data
- formats (.obj, .off), VRML and OpenGL
- no information about additional connectivity
- traversal is expensive

Edge-based data structures (winged-edge)



Vertex

coinciding edge
coordinates

Faces

coinciding edges

Edges

coinciding vertices

coinciding faces

prev. and next edge
of the left traverse

prev. and next edge
of the right traverse

- connectivity via edges
- clockwise ordering
- fast traversal through edges, vertices and faces
- case distinctions

Halfedge-based data structure

Vertex

halfedge pointing to the vertex
coordinates

Faces

coinciding halfedge

Edges

vertex it points to
adjacent face
next halfedge
previous halfedge
opposite halfedge

- edge is split into two halfedges
- consistent counter-clockwise ordering
- only for 2-manifolds

- C++ library – **OpenMesh**

